# Design of an 8-nsec-search 72-bit-word Content-Addressable Memory Using Phase-Change Devices 

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## 1. Introduction

Storage density and reliability of content-addressable memories (CAMs) have been continuously improved from the viewpoints of both device technologies and architecture [1]-[5]. Although CAMs for network routers require higher levels of reliability, errors in their stored data cannot be corrected during search operations. It is thus essential to improve the reliability of storage devices.

Phase-change devices not only have good scalability and resistance ratios higher than ten but also have good robustness against ionizing radiation [6]. These advantages make phase-change devices potentially suitable for a highly reliable CAM. Accordingly, in this study, we designed a CAM using phase-change devices and clarified its architecture, performance, and required resistances.

## 2. Architecture

The proposed resistive CAM cell incorporates two storage devices, SD0 and SD1 (i.e., phase-change devices) in an XNOR circuit (Fig. 1). Although this simple structure suppresses the area of a memory cell, it requires a careful circuit design to discriminate match-line current, which includes leakage current in one half of the memory cell (that is, a memory element) whose search line is activated (Table 1).

In the case of a conventional plane match-line structure, the ratio of match-line currents under the unmatched condition to those under the matched condition, $\alpha_{\mathrm{P}}$, is expressed as follows:

$$
\begin{equation*}
\left\{I_{O N}+(n-1) \times I_{L}\right\} /\left(n \times I_{L}\right) \leq \alpha_{P} \leq I_{O N} / I_{L} . \tag{1}
\end{equation*}
$$

Here, $n, I_{\mathrm{ON}}$, and $I_{\mathrm{L}}$ are the number of memory cells located on the match line, the turn-on current under the unmatched condition, and the leakage current under the matched condition, respectively. Since 72 memory cells are located on the match line in a standard CAM, leakage currents cannot be neglected. For example, when a match-line sensing circuit requires $\alpha_{\mathrm{P}} \geq 10, I_{\mathrm{ON}}$ must be more than 1,000 times larger than $I_{\mathrm{L}}$. This means that the phase-change resistance under amorphous state, $R_{0}$, must be more than 1,000 times higher than the resistance under crystalline state, $R_{1}$. This restriction complicates the design of a device.

To reduce resistance ratio $\beta_{\mathrm{R}}\left(=R_{0} / R_{1}\right)$, a one-hot-spot block code (OB code) (Fig. 2) and a hierarchical match-line structure (Fig. 5(a)) [2] are implemented in the proposed resistive CAM. In the case of per-two-bit coding, the ratio of match-line currents, $\alpha_{\mathrm{H}}$, is given by

$$
\begin{equation*}
\alpha_{H} \geq I_{O N} /\left(3 \times I_{L}\right) \tag{2}
\end{equation*}
$$

## 3. Performance of search operation

## Retention time and operating condition

The retention time of the resistive CAM during search operations, $t$, depends on the search cycle time and the search power consumed in the memory element, $P$. Although a higher pre-charge voltage, $V_{\mathrm{PC}}$, is desirable to perform a single-end sensing (Fig. 5(a)), it enhances destruction of stored data. Since look-up tables in network routers are, however, updated periodically, one-day retention is sufficient for CAMs targeting this usage. In this case, the operating condition, namely, 8 -ns cycle time (including $1.5-\mathrm{ns}$ search-line activation) and $0.85-\mathrm{V}$ pre-charge, is derived from Equation (3) [7] (Fig. 3).

$$
\begin{equation*}
t=D \exp \left[E /\left\{k T_{a}+\frac{P}{P_{1}}\left(\frac{E}{\ln \left(t_{1} / D\right)}-k T_{a}\right)\right\}\right] \tag{3}
\end{equation*}
$$

## Memory element

The area of the memory element, based on $130-\mathrm{nm}$ CMOS technology, was evaluated. In this evaluation, the effective area, including the sensing transistor of the match line and a pre-charge switch, was $200 F^{2}$ ( $F$ : minimum feature size), as listed in Table 2. Moreover, typical models of phase-change resistances for circuit simulations have the same non-linearity as those of Ref. [8] (Fig. 4).
Performance evaluation and required resistances
In a circuit simulation under eight operating conditions, 8 -ns cycle search operations were performed successfully (Figs. 5 and 6). As shown in Figure 5(b), 0.9-ns search-line activation is short enough for one-day retention. As shown in Figure 6, for the match line at $0.85 \mathrm{~V}, R_{0}$ must be higher than $12 \mathrm{M} \Omega$ and $R_{1}$ must be lower than $97 \mathrm{k} \Omega$. The minimum of the resistance ratio, $\beta_{\mathrm{R}}$, is 123 and the minimum ratio of match-line currents, $\alpha_{\mathrm{H}}$, is 40 .

## 4. Summary

A resistive content-addressable memory (CAM) using phase-change devices was designed. To reduce resistance ratio, $\beta_{\mathrm{R}}\left(=R_{0} / R_{1}\right)$, of the CAM, a hierarchical match-line structure and an OB code were implemented. As for the phase-change devices, for 8 -ns search operation, the following characteristics are required: one-day search retention, $R_{0}$ of more than $12 \mathrm{M} \Omega$ and $R_{1}$ of less than $97 \mathrm{k} \Omega$ (i.e., $\beta_{\mathrm{R}}=123$ ).

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(a) Conventional

(b) Proposed Fig. 1: Concept of memory cell


Voltage on match line, $V_{\mathrm{PC}}(\mathrm{V})$
Fig. 3: Calculated search-retention time


Fig. 4: I-V characteristics of phase-change devices

Table 1: Comparison of match-line currents during search operation

|  | Plane match line with $n$ cells | Hierarchical match line featuring OB code |
| :---: | :---: | :---: |
| Matched | (Stored data ' 0 ') |  |
| current | $n \times I_{L}$ | $I_{o N} \sim 4 \times I_{o v}$ |
| Unmatched |  |  |
| current | $I_{o \mathrm{ov}}+(n-1) \times I_{L} \sim n \times I_{\text {OV }}$ | $0 \sim 3 \times I_{L}$ |


| $\begin{gathered} \text { Data } \\ \text { combination } \end{gathered}$ | Conventional ternary code |  | One-hot-spot block code |
| :---: | :---: | :---: | :---: |
|  | $\begin{array}{\|l} \text { Notation } \\ \left(a_{1}, a_{0}\right) \end{array}$ | Stored-data pattern $\left(N_{\mathrm{a} 1}, N_{\mathrm{b} 1} \_N_{\mathrm{a} 0}, N_{\mathrm{bo}}\right)$ | Stored-data pattern $\left(N_{3}, N_{2}, N_{1}, N_{0}\right)$ |
| 0 | 00 | 01_01 | 0001 |
| 1 | 01 | 01_10 | 0010 |
| 2 | 10 | 10_01 | 0100 |
| 3 | 11 | 10_10 | 1000 |
| 0,1 | 0x | 01_00 | 0011 |
| 2, 3 | 1X | 10_00 | 1100 |
| 0,2 | x0 | 00_01 | 0101 |
| 1, 3 | X1 | 00_10 | 1010 |
| 1,2 | - | - | 0110 |
| 0,3 | - | - | 1001 |
| 0-2 | - | - | 0111 |
| 1-3 | - | - | 1110 |
| 0, 2, 3 | - | - | 1101 |
| 0, 1, 3 | - | - | 1011 |
| 0-3 | xX | 00_00 | 1111 |
| None | - | - | 0000 |

(a) Code mapping

(b) Conventional word structure (C) Hierarchical word structure Fig. 2: Concept of one-hot-spot block code (OB code) Table 2: Comparison of memory-cell areas

| SRAM-based [9] | DRAM-based [9] | PCM-based (effective) |
| :---: | :---: | :---: |
| $7.63 \mu \mathrm{~m}^{2}$ | $3.59 \mu \mathrm{~m}^{2}$ | $3.29 \mu \mathrm{~m}^{2}\left(6.77 \mu \mathrm{~m}^{2}\right)$ |
| $451 \mathrm{~F}^{2}$ | $212 \mathrm{~F}^{2}$ | $194 F^{2}\left(400 F^{2}\right)$ |


(a) Schematic of memory array

|  | $1^{\text {st }}$ cycle | $2^{\text {nd }}$ cycle |
| :--- | :--- | :--- |
| Stored data | $1000 \_1000 \_\ldots \_\mathbf{1 0 0 0}$ | $1000 \_1000 \_\ldots \_1000$ |
| Search key (SLs) | $1000 \_1000 \_\ldots \_\mathbf{1 0 1 1}$ | $1000 \_1000 \_\ldots \_0111$ |
| Match line, ML | discharged (ION+2xIL) | mid-level (3xIL: worst) |
| Operation | matched | unmatched |


(b) Simulated waveforms

Fig. 5: Search operations


| Condition |  |  | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Voltage | VDD | V | 1.35 | 1.35 | 1.35 | 1.35 | 1.65 | 1.65 | 1.65 | 1.65 |
|  | VPC | $\checkmark$ | 0.65 | 0.65 | 0.65 | 0.65 | 0.75 | 0.75 | 0.75 | 0.75 |
| Temperature |  | c | 125 | 125 | -40 | -40 | 125 | 125 | -40 | -40 |
| S |  |  |  | fast |  | fas |  |  |  |  |

Fig. 6: Required resistances for 8-ns-cycle search operation

